

Type Theorists HATE Him!

Learn this **ONE WEIRD TRICK** to fake dependent types in a language that doesn't support them

LEARN THE TRUTH NOW

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Dependent types



Dependent types... in Idris

```
mkSingle : (x : Bool) ->  
          if x then Nat else List Nat  
mkSingle True = 0  
mkSingle False = []
```



Dependent types... in Haskell?

```
mkSingle :: (x :: Bool) ->
           if x then Nat else [Nat]
mkSingle True = 0
mkSingle False = []
```

Dependent types... in Haskell?

```
mkSingle :: (x :: Bool) ->
           if x then Nat else [Nat]
mkSingle True = 0
mkSingle False = []
```

```
<interactive>:1:28: error: parse error on input 'if'
```

:(



There's hope yet

- λ Haskell *can* support dependent types, if you're willing to squint
- λ We'll need to enable a modest number of GHC extensions



There's hope yet

λ Haskell *can* support dependent types, if you're willing to squint

λ We'll need to enable a modest number of GHC extensions

```
{-# LANGUAGE DefaultSignatures #-}
{-# LANGUAGE EmptyCase #-}
{-# LANGUAGE ExistentialQuantification #-}
{-# LANGUAGE FlexibleContexts #-}
{-# LANGUAGE FlexibleInstances #-}
{-# LANGUAGE GADTs #-}
{-# LANGUAGE InstanceSigs #-}
{-# LANGUAGE KindSignatures #-}
{-# LANGUAGE RankNTypes #-}
{-# LANGUAGE ScopedTypeVariables #-}
{-# LANGUAGE TemplateHaskell #-}
{-# LANGUAGE TypeFamilies #-}
{-# LANGUAGE TypeInType #-}
{-# LANGUAGE TypeOperators #-}
{-# LANGUAGE UndecidableInstances #-}
```



Singleton types

```
data family Sing :: k -> Type
```




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```
data instance Sing :: Bool -> Type where  
  SFalse :: Sing False  
  STrue  :: Sing True
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Singleton types

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data family Sing :: k -> Type
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```
data instance Sing :: Bool -> Type where
  SFalse :: Sing False
  STRue  :: Sing True
```

```
data instance Sing (z :: Bool)
  = (z ~ False) => SFalse
  | (z ~ True)  => STRue
```

Singleton types

```
data family Sing :: k -> Type
```

```
data instance Sing :: Bool -> Type where
  SFalse :: Sing False
  STrue  :: Sing True
```

```
data instance Sing (z :: Bool)
  = (z ~ False) => SFalse
  | (z ~ True)  => STrue
```



Singleton types

```
data family Sing :: k -> Type
```

```
data Nat = Z | S Nat
```

```
data instance Sing :: Nat -> Type where  
  SZ :: Sing Z  
  SS :: Sing (n :: Nat) -> Sing (S n)
```

Singleton types

```
data family Sing :: k -> Type
```

```
data Nat = Z | S Nat
```

```
data instance Sing :: Nat -> Type where
  SZ :: Sing Z
  SS :: Sing (n :: Nat) -> Sing (S n)
```

```
data instance Sing (z :: Nat)
  = (z ~ Z) => SZ
  | forall! (n :: Nat). (z ~ S n) => SS n
```



Dependent types... in Idris (redux)

```
mkSingle : (x : Bool) ->
  if x then Nat else List Nat
mkSingle True = 0
mkSingle False = []
```




Dependent types... in Haskell? (Heck yeah!)

```
type family
If (c :: Bool) (t :: k) (f :: k) :: k where
If True  t f = t
If False t f = f
```



Dependent types... in Haskell? (Heck yeah!)

```
type family
If (c :: Bool) (t :: k) (f :: k) :: k where
If True  t f = t
If False t f = f

mkSingle :: Sing (x :: Bool) ->
          If x Nat [Nat]
mkSingle STRue  = ⊙
mkSingle SFalse = []
```



Dependent types... in Haskell?
(Heck yeah!)

type $\Pi = \text{Sing}$



Dependent types... in Haskell? (Heck yeah!)

type $\Pi = \text{Sing}$

```
mkSingle ::  $\Pi$  (x :: Bool) ->
           If x Nat [Nat]
mkSingle STRue =  $\odot$ 
mkSingle SFaLse = []
```



What else can we fake emulate?

λ Dependent pattern matching



What else can we fake emulate?

λ Dependent pattern matching... in Idris

```
data Vect : Nat -> Type -> Type where
  Nil  : Vect Z a
  Cons : a -> Vect n a -> Vect (S n) a
```



What else can we fake emulate?

λ Dependent pattern matching... in Idris

```
data Vect : Nat -> Type -> Type where
  Nil  : Vect Z a
  Cons : a -> Vect n a -> Vect (S n) a
```

```
Len : Vect n a -> Nat
Len Nil = 0
Len (Cons x xs) = 1 + Len xs
```



What else can we fake emulate?

λ Dependent pattern matching... in Idris

```
data Vect : Nat -> Type -> Type where
  Nil  : Vect Z a
  Cons : a -> Vect n a -> Vect (S n) a
```

```
Len : {n : Nat} -> Vect n a -> Nat
Len {n=Z} Nil = 0
Len {n=S k} (Cons x xs) = 1 + Len xs
```




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λ Dependent pattern matching... in Haskell!



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λ Dependent pattern matching... in Haskell!

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data Vect :: Nat -> Type -> Type where
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What else can we fake emulate?

λ Dependent pattern matching... in Haskell!

```
data Vect :: Nat -> Type -> Type where
  Nil  :: Vect Z a
  Cons :: a -> Vect n a -> Vect (S n) a
```

```
Len :: Sing (n :: Nat) -> Vect n a -> Nat
Len SZ Nil           = 0
Len (SS k) (Cons x xs) = 1 + Len xs
```



What else can we fake emulate?

λ Dependent pattern matching... in Haskell!

```
class SingI (a :: k) where
  sing :: Sing (a :: k)
```

```
instance SingI Z where
  sing = SZ
```

```
instance SingI n => SingI (S n) where
  sing = SS sing
```



What else can we fake emulate?

λ Dependent pattern matching... in Haskell!

```
Len :: Sing (n :: Nat) -> Vect n a -> Nat
Len SZ Nil           = 0
Len (SS k) (Cons x xs) = 1 + Len xs
```



What else can we fake emulate?

λ Dependent pattern matching... in Haskell!

```
Len :: Sing (n :: Nat) -> Vect n a -> Nat
Len SZ Nil = 0
Len (SS k) (Cons x xs) = 1 + Len xs

Len' :: forall1 (n :: Nat). SingI n =>
      Vect n a -> Nat
Len' = Len (sing :: n)
```

- λ With enough elbow grease, one can simulate a great deal of dependently typed code
- λ Impress your friends at the bar! Be the envy of your family!
- λ <http://hackage.haskell.org/package/singleton>

Any questions?